



BM 307 Dosya Organizasyonu (File Organization)

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Bilgisayar Mühendisliği Bölümü



Konular

- B-Tree and Derivatives
 - B-Trees
 - B#-Trees
 - B+-Trees
 - Değerlendirme



B-Tree and Derivatives

B-Trees

- B-Tree sequential ve direct erişimde iyi performansa sahiptir.
- Binary tree'lerde branching factor ikiden büyük olamaz. B-Tree'lerde teorik olarak limit yoktur.
- B-Tree'ler multiway search tree'dir ve bir node'da çok sayıda kayıt tutulabilir
- B-Tree'ler diğer ağaçlardan farklı olarak bottom-up yapılandırılırlar.
- Herhangi bir döndürmeye gerek kalmadan (AVL tree ve IPR tree'deki gibi) otomatik olarak balance edilirler.



B-Tree and Derivatives

B-Trees

- B-tree aşağıdaki kriterlere göre oluşturulur:

1- $d \leq \text{keys} \leq 2d$ (root node için $1 \leq \text{keys} \leq 2d$)

$d+1 \leq \text{keys} \leq 2d+1$ (root node için $2 \leq \text{pointers} \leq 2d+1$)

2- Bütün yaprak node'lar aynı seviyededir.

Pointer	Key (record)	Pointer	Key (record)	Pointer
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B-Tree and Derivatives

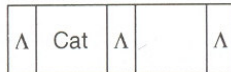
B-Trees

Örnek

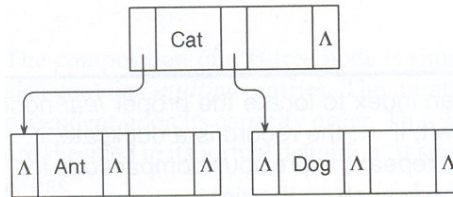
$d = 1$ (capacity order)

cat, ant, dog, cow, rat, pig, gnu

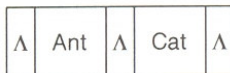
Cat insert edildi



Dog insert edildi



Ant insert edildi



B-Tree and Derivatives

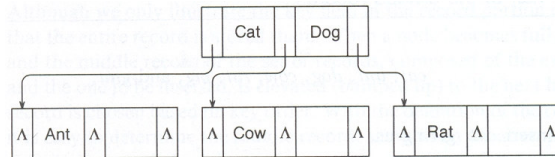
B-Trees

Örnek

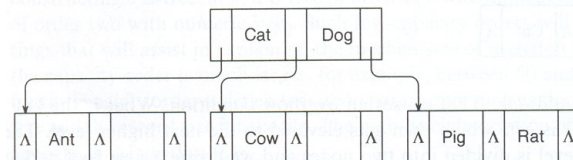
$d = 1$ (capacity order)

cat, ant, dog, cow, rat, pig, gnu

Cow ve Rat insert edildi



Pig insert edildi





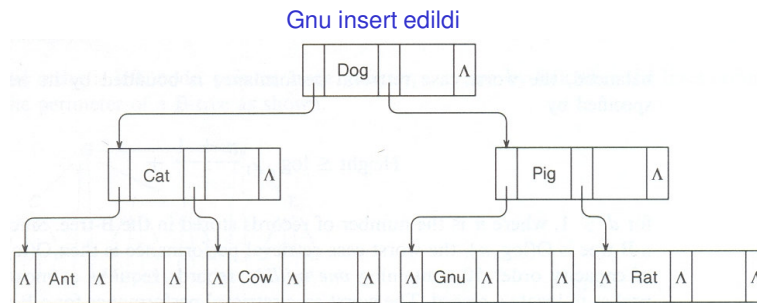
B-Tree and Derivatives

B-Trees

Örnek

$d = 1$ (capacity order)

cat, ant, dog, cow, rat, pig, gnu



Packing factor = depolanan kayıt sayısı / kullanılan yer sayısı
= 7 / 14 = 50 %



B-Tree and Derivatives

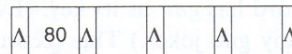
B-Trees

Örnek

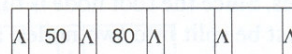
$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150

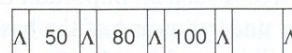
1. Insert 80.



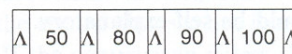
2. Insert 50.



3. Insert 100.



4. Insert 90





B-Tree and Derivatives

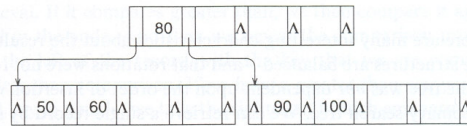
B-Trees

■ Örnek

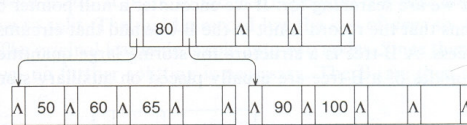
$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150

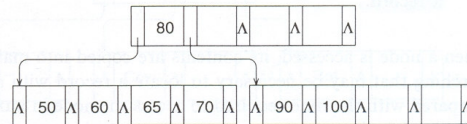
5. Insert 60.



6. Insert 65.



7. Insert 70.



B-Tree and Derivatives

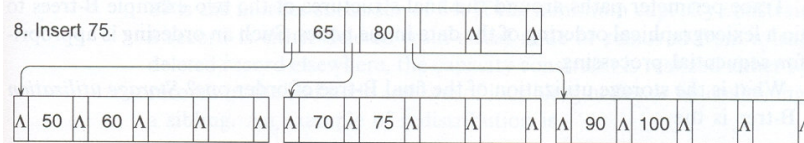
B-Trees

■ Örnek

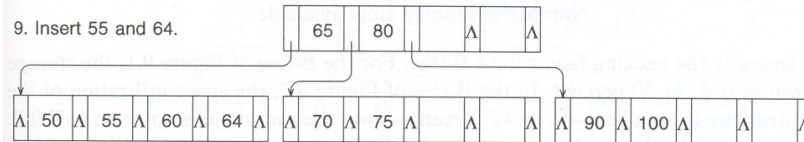
$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150

8. Insert 75.



9. Insert 55 and 64.





B-Tree and Derivatives

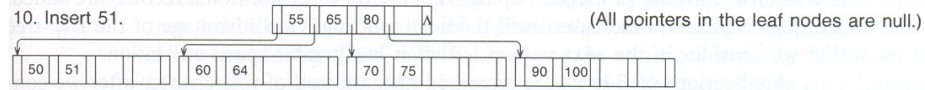
B-Trees

Örnek

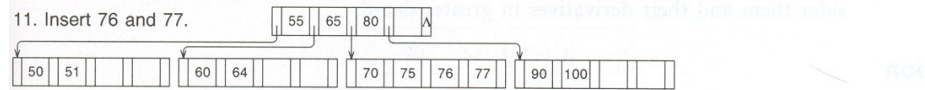
$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150

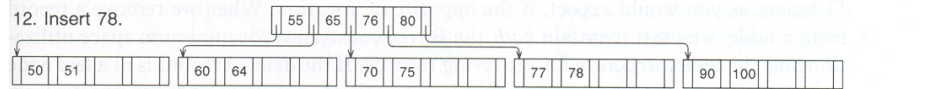
10. Insert 51. (All pointers in the leaf nodes are null.)



11. Insert 76 and 77.



12. Insert 78.



B-Tree and Derivatives

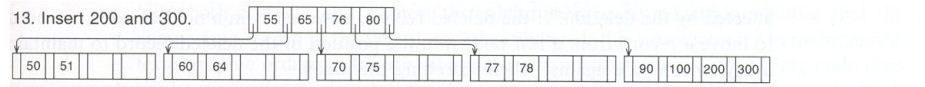
B-Trees

Örnek

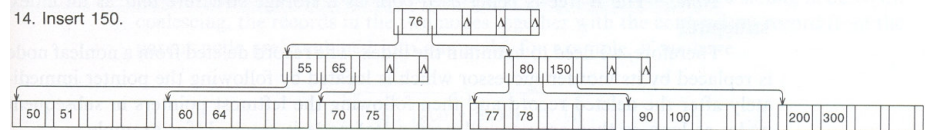
$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150

13. Insert 200 and 300.



14. Insert 150.



Packing factor = depolanan kayıt sayısı / kullanılan yer sayısı
= 17 / 36 = 47 %



B-Tree and Derivatives

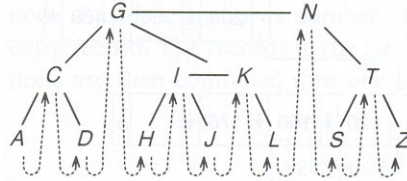
B-Trees

Değerlendirme

- B-Tree'lerde retrieval probe kayda ulaşma değil node'a ulaşmaktır
- Ağacın yüksekliği aşağıdaki gibi ifade edilir

$$\text{Height} \leq \log_{d+1} \frac{n+1}{2} + 1$$

- Worst case retrieval performansı $O(\log_d n)$ 'dir. n, B-tree'deki kayıt sayısıdır
- Capacity order'ı 50 olan ve 1 milyon kayıt bulunduran B-tree'de bir kayda ulaşmak için en fazla 4 retrieval probe gerekir.
- B-tree'deki kayıtlara sıralı ulaşmak için inorder dolaşma yapılır.

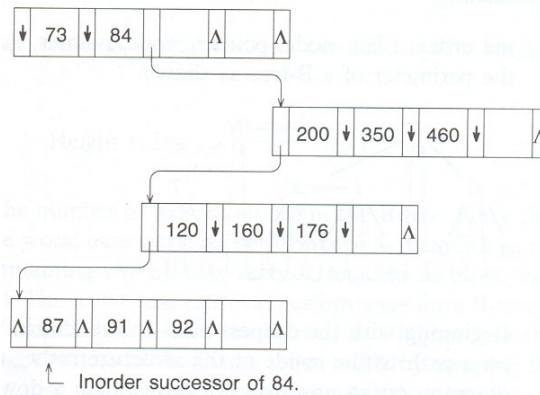


B-Tree and Derivatives

B-Trees

Silme

- Bir nonleaf node üzerinden kayıt silindiğinde inorder takipçisi yerine yazılır



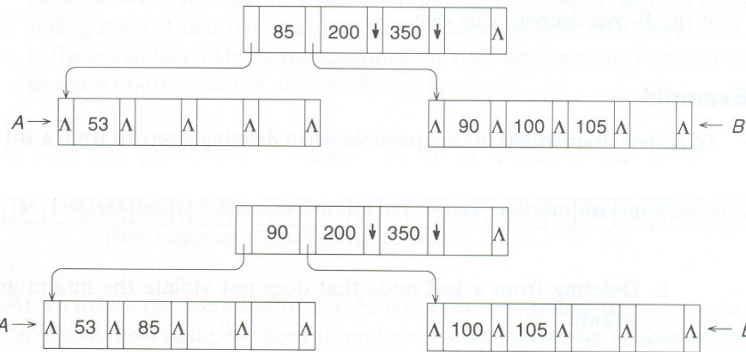


B-Tree and Derivatives

B-Trees

- Silme

- Bir node'daki kayıt sayısı minimum kapasite'den aşağı düşerse ve kardeş node'u fazla kayda sahipse, parent ve kardeş node ile yeniden düzenleme yapılır

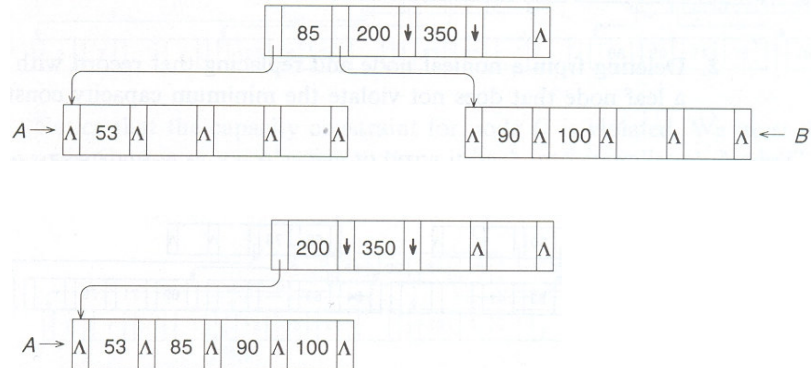


B-Tree and Derivatives

B-Trees

- Silme

- İki kardeş node minimum kapasitenin altına düşerse ikisi ve parent node'daki kayıt birleştirilir

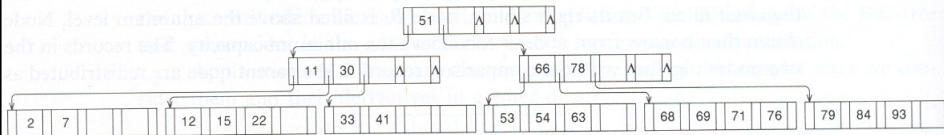




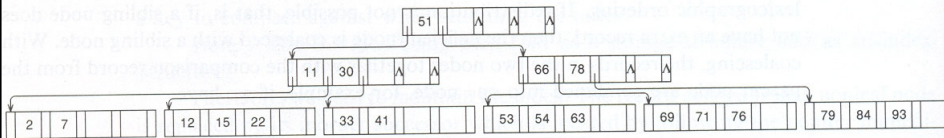
B-Tree and Derivatives

B-Trees

- Silme - Örnek
 - Minimum kapasitenin üstündeki yapraktan kayıt silinmesi



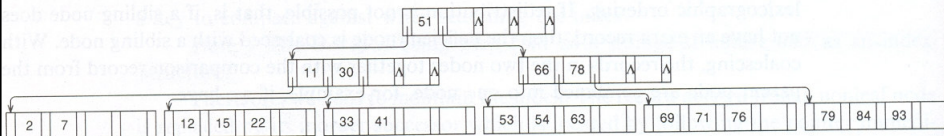
68 silindi



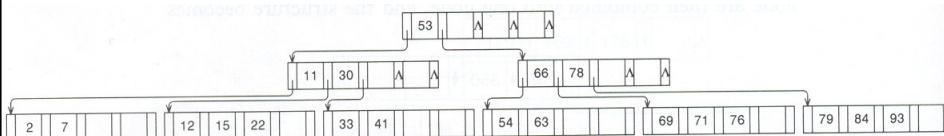
B-Tree and Derivatives

B-Trees

- Silme - Örnek
 - Bir nonleaf node'da kayıt silinmesi ve minimum kapasitenin üzerindeki bir node'dan kayıt aktarılması



51 silindi



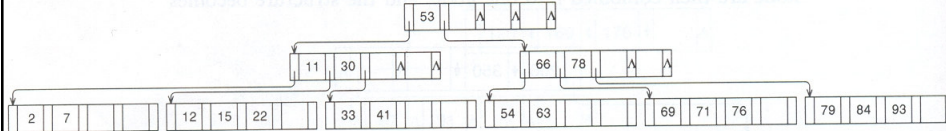


B-Tree and Derivatives

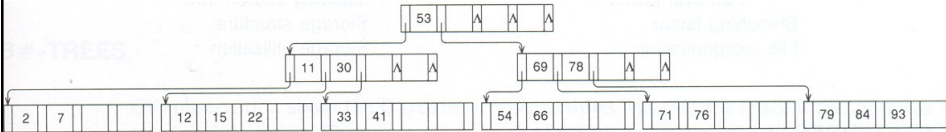
B-Trees

■ Silme - Örnek

- Bir leaf node'da kayıt silinmesi ve minimum kapasitenin altına düşülmesi



63 silindi

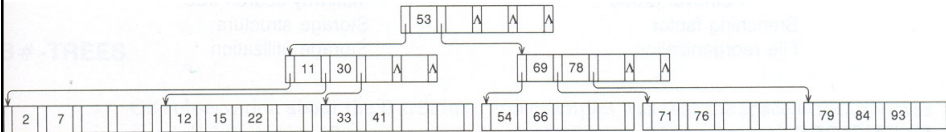


B-Tree and Derivatives

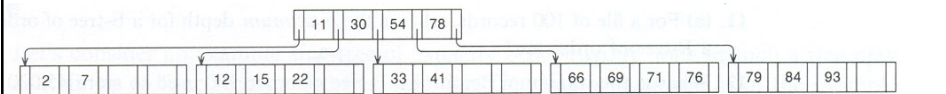
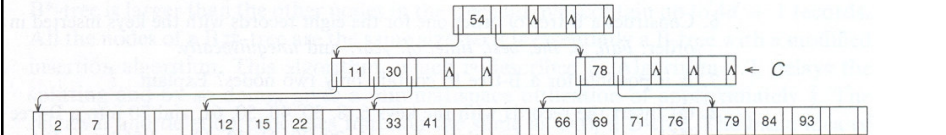
B-Trees

■ Silme - Örnek

- Bir leaf node'da kayıt silinmesi ve minimum kapasitenin altına düşülmesi ve node'ların birleştirilmesi



53 silindi





B-Tree and Derivatives

B#-Trees

- B-tree'lerde bir node dolunca bölme işlemi yapılmaktadır
- Bölme sonucunda oluşan iki node'da yarı yarıya doludur
- B#-tree'lerde bölme işlemi geciktirilerek node'ların doluluk oranı artırılır
- Ortalama Insert süresi uzar ve ağacın yüksekliği daha azdır.
- Tüm ağacın packing faktor değeri B-tree'lere göre daha yüksektir
- B# - tree'lerde retrieval performansı daha yüksektir
- Literatürde B*-tree şeklinde variantları vardır



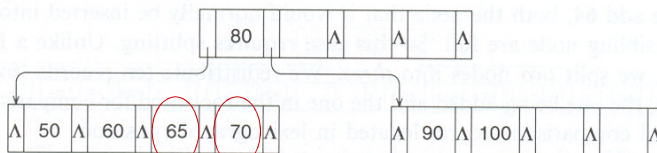
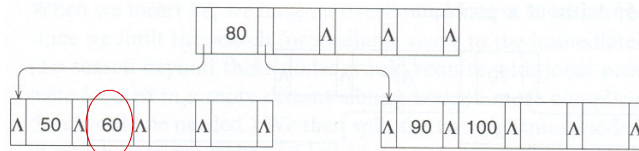
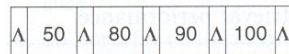
B-Tree and Derivatives

B#-Trees

- Örnek

$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150





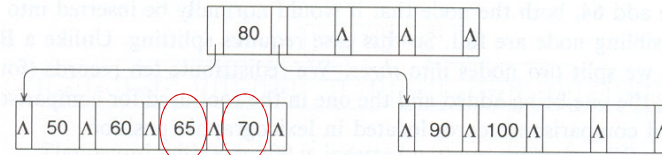
B-Tree and Derivatives

B#-Trees

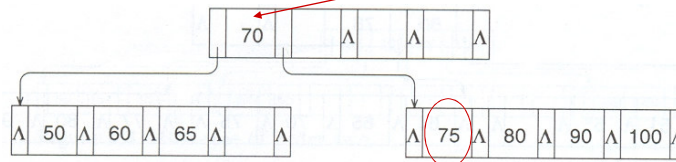
■ Örnek

$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150



comparison.record = $\lfloor (r + 1) / n \rfloor$ r, dağıtılacak toplam kayıt sayısı
 $\lfloor 9/2 \rfloor = 4$



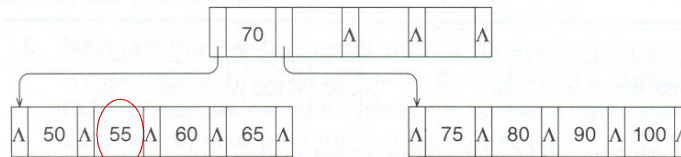
B-Tree and Derivatives

B#-Trees

■ Örnek

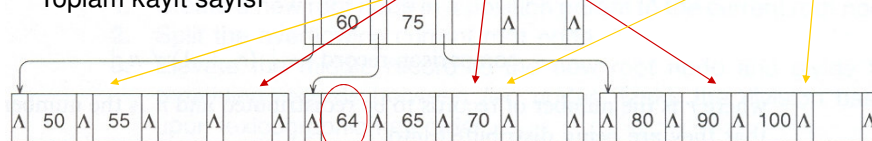
$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150



$\lfloor (10 - 2) / 3 \rfloor = 2$ $\lfloor 6 / 2 \rfloor = 3$

Toplam kayıt sayısı





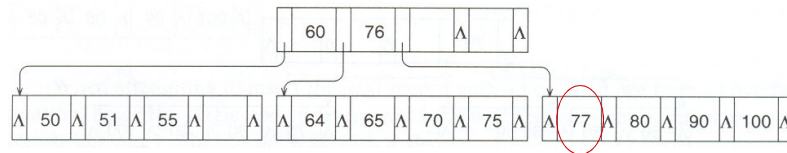
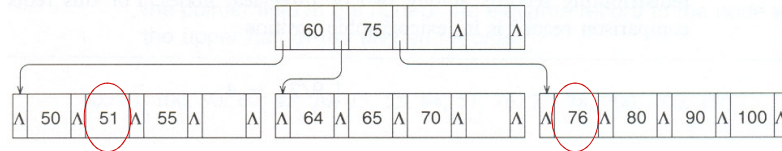
B-Tree and Derivatives

B#-Trees

■ Örnek

$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150



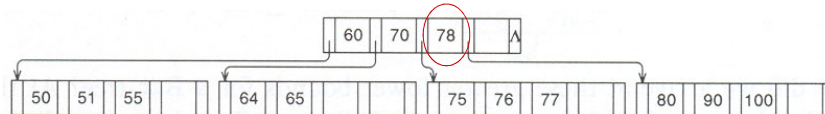
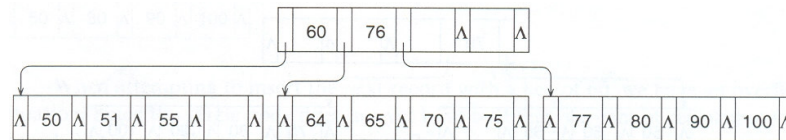
B-Tree and Derivatives

B#-Trees

■ Örnek

$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150



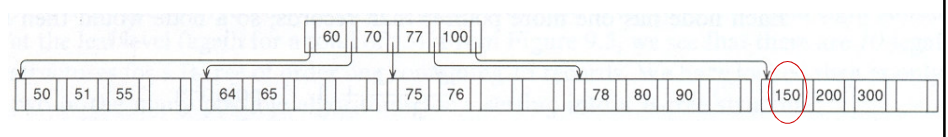
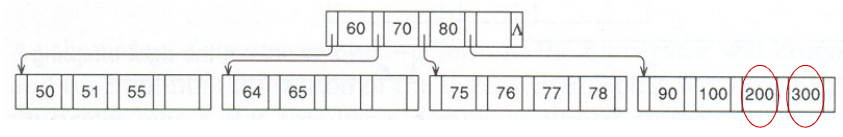
B-Tree and Derivatives

B#-Trees

Örnek

$d = 2$ (capacity order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150



B-Tree and Derivatives

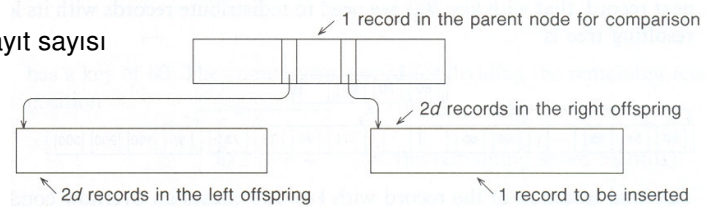
B#-Trees

Tanım

$$\frac{4d - 2}{3} \leq \text{keys} \leq 2d$$

$$\frac{4d + 1}{3} \leq \text{pointers} \leq 2d + 1$$

Toplam kayıt sayısı
 $4d + 2$



Node'lara dağıtılacak en az kayıt sayısı

$$\left\lceil \frac{4d}{3} \right\rceil$$

Minimum doluluk oranı

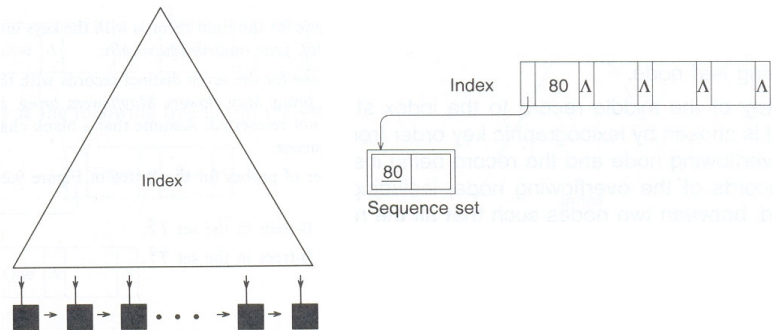
$$\frac{(4d - 2)/3}{2d} \rightarrow \frac{4d - 2}{6d} \rightarrow 2/3$$



B-Tree and Derivatives

B+ -Trees

- B+ -tree'lerde sıralı okuma için parent'a ulaşmaya gerek yoktur
- Bilgiler yapraklarda bulunur. Nonleaf node'lar sadece indeks için kullanılır
- Tüm yapraklar tek bağlı listeyle bağlanır
- B+ - tree'lerde indeks kısmı B-tree ile aynıdır. İndeks kısmındaki tüm kısıtlamalar ve işlemler B-tree ile aynıdır



B-Tree and Derivatives

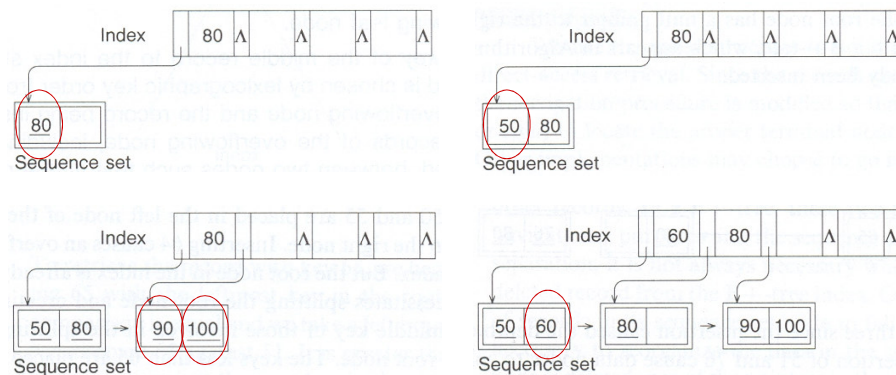
B+ -Trees

■ Örnek

$d = 2$ (capacity order)

$s = 1$ (sequence set order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150





B-Tree and Derivatives

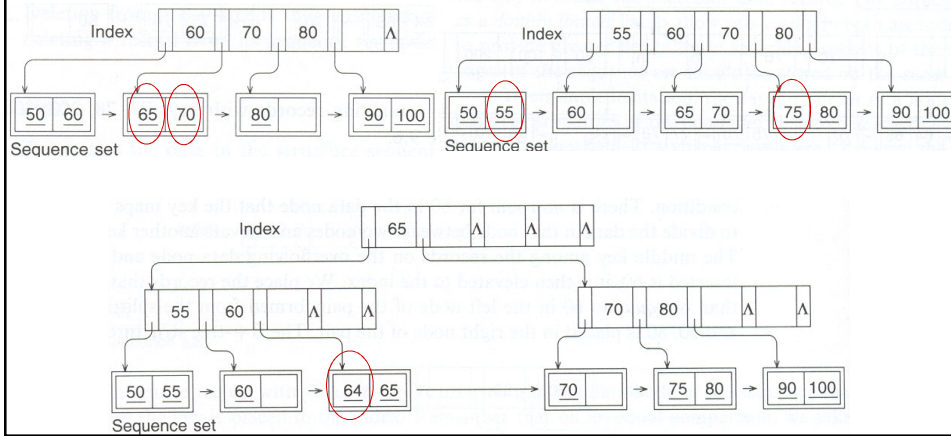
B+ -Trees

Örnek

$d = 2$ (capacity order)

$s = 1$ (sequence set order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150



B-Tree and Derivatives

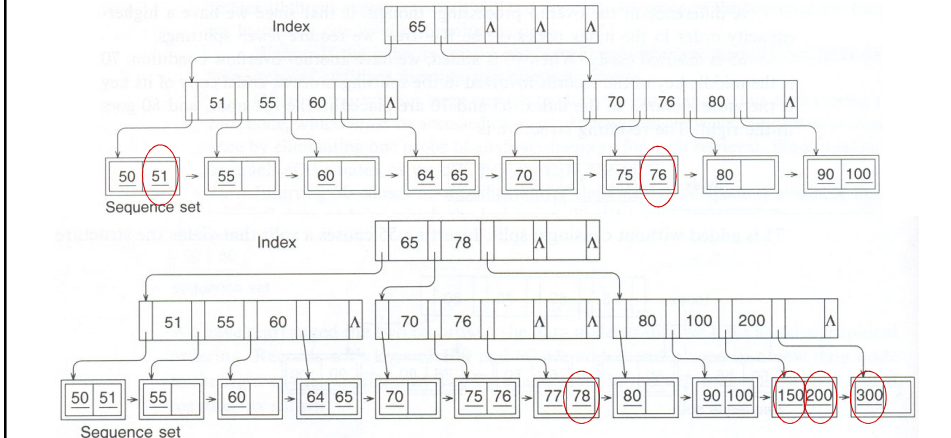
B+ -Trees

Örnek

$d = 2$ (capacity order)

$s = 1$ (sequence set order)

80,50,100,90,60,65,70,75,55,64,51,76,77,78,200,300,150





Binary Tree Structures

Haftalık Ödev

B+ - tree'lerde ve B# - tree'lerde silme işleminin nasıl yapıldığını araştırınız ve bir rapor hazırlayınız .